## Lecture 13

## **Arithmetic Functions**

Today - Arithmetic functions, the Möbius function

**(Definition) Arithmetic Function:** An **arithmetic function** is a function  $f: \mathbb{N} \to \mathbb{C}$ 

Eg.

 $\pi(n) = \text{ the number of primes } \leq n$ 

d(n) = the number of positive divisors of n

 $\sigma(n) =$  the sum of the positive divisors of n

 $\sigma_k(n) =$  the sum of the kth powers of n

 $\omega(n)$  = the number of distinct primes dividing n

 $\Omega(n) =$ the number of primes dividing n counted with multiplicity

Eg.

$$\sigma(1) = 1$$

$$\sigma(2) = 1 + 2 = 3$$

$$\sigma(3) = 1 + 3 = 4$$

$$\sigma(6) = 1 + 2 + 3 + 6 = 12$$

(**Definition**) **Perfect Number:** A **perfect number** n is one for which  $\sigma(n) = 2n$  (eg., 6, 28, 496, etc.)

Big open conjecture: Every perfect number is even.

**Note:** One can show that if n is an even perfect number, then  $n=2^{m-1}(2^m-1)$  where  $2^m-1$  is a Mersenne prime (Euler)

**(Definition) Multiplicative:** If f is an arithmetic function such that whenever (m,n)=1 then f(mn)=f(m)f(n), we say f is **multiplicative**. If f satisfies the stronger property that f(mn)=f(m)f(n) for all m,n (even if not coprime), we say f is **completely multiplicative** 

Eg.

$$f(n) = \begin{cases} 1 & n = 1 \\ 0 & n < 1 \end{cases}$$

is completely multiplicative. It's sometimes called 1 (we'll see why soon).

**Eg.**  $f(n) = n^k$  for some fixed  $k \in \mathbb{N}$  is also completely multiplicative

**Eg.**  $\omega(n)$  is not multiplicative (adds, but  $2^{\omega(n)}$  is multiplicative)

**Eg.**  $\phi(n)$  is multiplicative (by CRT)

**Note:** If f is a multiplicative function, then to know f(n) for all n, it suffices to know f(n) for prime powers n. This is why we wrote

$$\phi(p_1^{e_1} \dots p_r^{e_r}) = \prod p_i^{e_i - 1} (p_i - 1)$$

**(Definition) Convolution:** The **convolution** of two arithmetic functions f and g is f \* g defined by

$$(f*g)(n) = \sum_{d|n} f(d)g\left(\frac{n}{d}\right)$$

(summing over positive divisors of n). Compare this to convolution from calculus or differential equations

$$(f*g)(x)=\int_{t=-\infty}^{\infty}f(t)g(x-t)dt$$
 or 
$$\int_{t=0}^{x}f(t)g(x-t)dt \quad \text{ if } f(y),g(y) \text{ are } 0 \text{ for } y<0$$

Eg.

$$f * 1 = 1 * f = f$$
 for every  $f$  (check 1)

1 is the identity for convolution

**Theorem 48.** If f and g are multiplicative then f \* g is multiplicative.

*Proof.* Suppose m and n are coprime. Then any divisor of mn is of the form

 $d_1, d_2$ , where  $d_1|m$  and  $d_2|n$ , uniquely. So we have

$$(f * g)(mn) = \sum_{d|mn} f(d)g\left(\frac{mn}{d}\right)$$

$$= \sum_{d_1|m} \sum_{d_2|n} f(d_1d_2)g\left(\frac{m}{d_1} \cdot \frac{n}{d_2}\right)$$

$$= \sum_{d_1|m} \sum_{d_2|n} \underbrace{f(d_1)f(d_2)}_{\text{since }(d_1,d_2)=1} \underbrace{g\left(\frac{m}{d_1}\right)g\left(\frac{n}{d_2}\right)}_{\text{since }(\frac{m}{d_1},\frac{n}{d_2})=1}$$

$$= \left(\sum_{d_1|m} f(d_1)g\left(\frac{m}{d_1}\right)\right) \left(\sum_{d_2|n} f(d_2)g\left(\frac{n}{d_2}\right)\right)$$

$$= (f * g)(m)(f * g)(n)$$

**Eg.** Let U(n) = 1 for all n. Then for any arithmetic function f, we have

$$(f * U)(n) = \sum_{d|n} f(d) \underbrace{U\left(\frac{n}{d}\right)}_{-1} = \sum_{d|n} f(d)$$

This is usually called F(n).

If f is multiplicative, then F is multiplicative, by theorem (since U is obviously completely multiplicative) (that's theorem 4.4 in the book). In particular, we compute U\*U

$$(U*U)(n) = \sum_{d|n} 1 \cdot 1 = \text{number of divisors of } n = d(n)$$

so d(n) is multiplicative.

For a prime power  $p^{\alpha}$ , the number of divisors is  $\alpha + 1$ . So

$$d(p_1^{e_1} \dots p_r^{e_r}) = \prod_{i=1}^r (e_i + 1)$$

For the function  $r_k(n) = n^k$ , we have

$$(r_k * U)(n) = \sum_{d|n} d^k = \sigma_k(n)$$

which is therefore multiplicative.

Since

$$\sigma_k(p^{\alpha}) = 1 + p^k + \dots + p^{k\alpha} = \frac{p^{k(\alpha+1)} - 1}{p^k - 1}$$

we get

$$\sigma_k\left(\prod p_i^{e_i}\right) = \prod_{i=1}^r \frac{p_i^{k(e_i+1)} - 1}{p_i^k - 1}$$

**Theorem 49.** For any positive integer n, we have

$$\sum_{d|n} \phi(n) = n$$

In other words  $\phi * U = r_1$ .

*Proof.* Since both sides are multiplicative, enough to show it for prime powers.

$$r_1(p^{\alpha}) = p^{\alpha}$$

$$\sum_{d|p^{\alpha}} \phi(d) = \phi(1) + \phi(p) + \dots + \phi(p^{\alpha})$$

$$= 1 + (p-1) + p(p-1) + \dots + p^{\alpha-1}(p-1)$$

$$= 1 + p - 1 + p^2 - p + p^3 - p^2 + \dots + p^{\alpha} - p^{\alpha-1}$$

$$= p^{\alpha}$$

**Eg.** What is  $r_k * r_k$ ?

$$(r_k * r_k)(n) = \sum_{d|n} d^k \left(\frac{n}{d}\right)^k = \sum_{d|n} n_k = n^k d(n)$$

Other multiplicative functions:  $(\overline{D})$  since  $(\frac{mn}{D}) = (\frac{m}{D})(\frac{n}{D})$ 

There's an interesting multiplicative function - let  $\tau(n)$  (Ramanujan's tau function) be the coefficient of  $q^n$  in  $q \prod_{i=1}^{\infty} (1-q^i)^{24}$ . Then

$$\begin{split} \tau(1) &= 1 \\ \tau(2) &= -24 \\ \tau(3) &= 252 \\ \tau(6) &= -6048 = -24 \cdot 252 = \tau(2)\tau(3) \\ \text{etc.} \end{split}$$

**Theorem 50.**  $\tau(n)$  is multiplicative (deep)

Proof uses modular forms.

A famous open conjecture is **Lehmer's conjecture**:  $\tau(n) \neq 0$  for every natural number n.

**Proposition 51.** f \* g = g \* f for any arithmetic functions f, g (ie., convolution is commutative)

Proof.

$$(f * g)(n) = \sum_{d|n} f(d)g\left(\frac{n}{d}\right)$$

as d ranges over divisors of n, so does  $\frac{n}{d} = d'$  so we write

$$(f*g)(n) = \sum_{d'|n} f(d')g(\frac{n}{d'}) = \sum_{d|n} f\left(\frac{n}{d}\right)g\left(\frac{n}{\frac{n}{d}}\right) = \sum_{d|n} f\left(\frac{n}{d}\right)g(d) = (g*f)(n)$$

Proposition 52.

$$f * (g * h) = (f * g) * h$$

(ie., \* is associative)

Proof left as exercise.

The Möbius mu function is defined to be

$$\mu(n) = \begin{cases} (-1)^{\omega(n)} & \text{if } n \text{ is square free} \\ 0 & \text{otherwise} \end{cases}$$

remembering that  $\omega(n)$  was additive, it's easy to see  $\mu(n)$  is a multiplicative function

$$\mu(n) = \begin{cases} (-1)^r & \text{if } n = p_1 \dots p_r \text{ for distinct primes } p_i \\ 0 & \text{if some } p^2 \text{ divides } n \end{cases}$$

Remember the function U(n) = 1 for all n.

Theorem 53.

$$\mu*U=\mathbb{1}\begin{cases} 1 & \textit{if } n=1\\ 0 & \textit{otherwise} \end{cases}$$

*Proof.*  $\mu, U$  are multiplicative, so  $\mu*U$  is too. So enough to show that  $(\mu*U)(1)=1$  and  $(\mu*U)(p^{\alpha})=0$  for prime powers  $p^{\alpha}$ .

The first is trivial:

$$(\mu * U)(1) = \sum_{d|1} \mu(1)U(1) = 1 \cdot 1 = 1$$
$$(\mu * U)(p^{\alpha}) = \sum_{d|p^{\alpha}} \mu(p^{\alpha})U(p^{\alpha})$$
$$= \mu(1) + \mu(p)$$
$$= 1 + (-1)$$
$$= 0$$

So 
$$\mu * U = 1$$
.

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